

MIFARE Classic: Completely Broken

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Introduction

- MIFARE Classic
 - Owned by NXP Semiconductors, Inc.
 - The most widely deployed RFID technology
 - Over 1 billion cards sold
 - Main uses
 - Public transportation ticketing systems
 - Access control systems
 - Reverse-engineered in late 2008 by European hackers
- In this talk, I will report our first-hand experience attacking a real MIFARE Classic system

Acknowledgments

- K. Nohl, D. Evans, and H. Plötz. “Reverse-engineering a cryptographic RFID tag.” In USENIX Security Symposium 2008
- F. D. Garcia, P. van Rossum, R. Verdult, and R. W. Schreur. “Wirelessly pickpocketing a MIFARE Classic card.” In IEEE Symposium on Security and Privacy 2009
- M.-Y. Chih, J.-R. Shih, B.-Y. Yang, J. Ding, and C.-M. Cheng. “MIFARE Classic: Practical attacks and defenses.” In CISC 2010

Outline

- Overview of MIFARE Classic
 - Memory layout
 - Communication protocol
 - Authentication protocol
 - CRYPTO-1 stream cipher
- Principal technique: known-plaintext attack
- Reader-based attacks
- Sniffer-based attacks
- Concluding remarks

Jargon of the Trade

- MIFARE Classic is based on the ISO/IEC 14443 Type A 13.56 MHz contactless smart card standard
 - A reader is referred to as a PCD (Proximity Coupling Device), whereas a card/tag, PICC (Proximity Integrated Circuit Card)
 - We will use these terms interchangeably with readers, cards, and tags

Memory Layout

Memory size	1 KB	4 KB
# Blocks	64	256
# Sectors	16	40
# Blocks in a sector	4	4 or 12
Example		

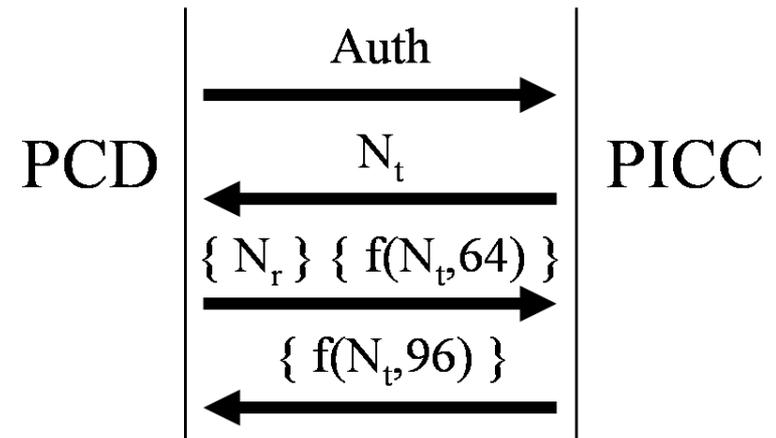
Sector number	Block number	Content (16 Bytes)							
0	0	UID, BCC, Manufacturer (Read Only)							
	1.Data/Value	Data or Value							
	2.Data/Value	Data or Value							
	3.Tail	Key A	Access cond.	U	Key B				
1	4.Data/Value	Data or Value							
	5.Data/Value	Data or Value							
	6.Data/Value	Data or Value							
	7.Tail	Key A	Access cond.	U	Key B				
⋮									
15	60.Data/Value	Value	$\overline{\text{Value}}$	Value	00	ff	00	ff	
	61.Data/Value	Value	$\overline{\text{Value}}$	Value	00	ff	00	ff	
	62.Data/Value	Data/Value							
	63.Tail	Key A	Access cond.	U	Key B				
MIFARE Classic 1K Memory Layout									

● Block:

- Data – 16 bytes
- Value – 4 bytes
- Sector tail – access control

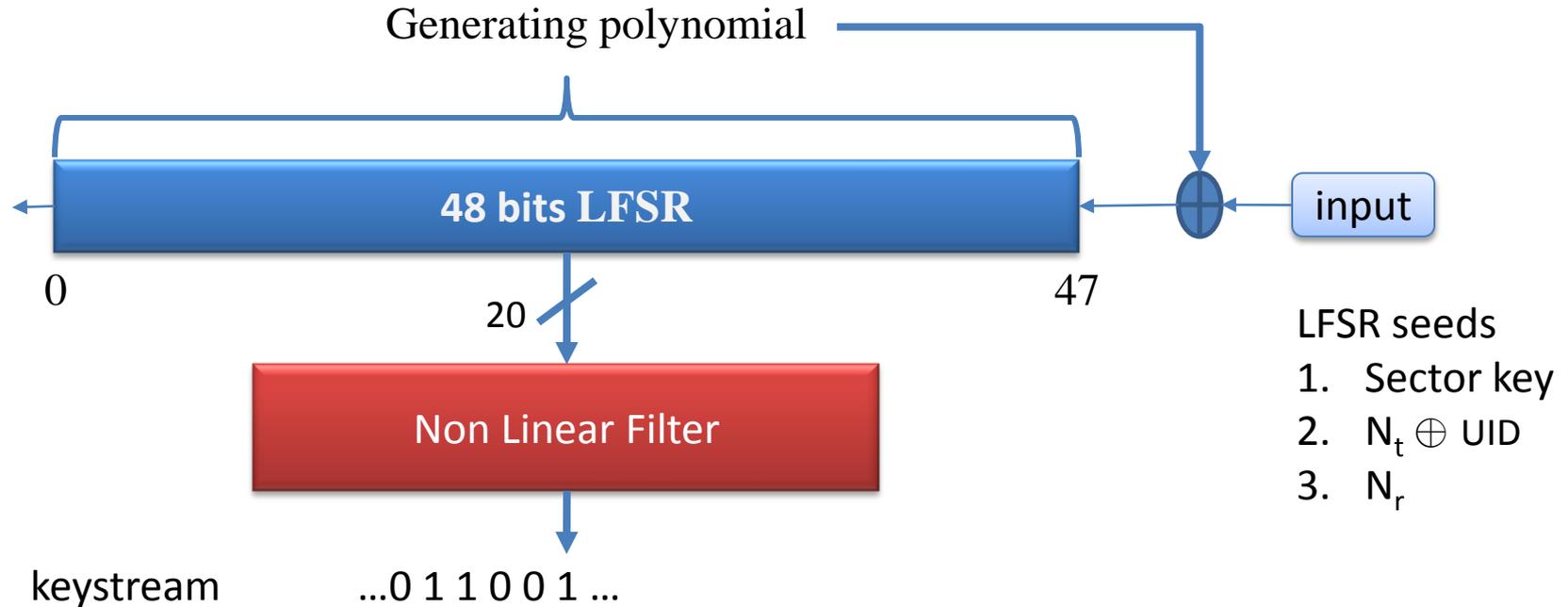
Communication and Authentication

1. Anti-collision (UID)
2. Authentication (key A/B)
3. Memory operations
 - ① Read
 - ② Write
 - ③ Increment, decrement, restore
 - ④ Halt



Cryptographic Primitive

The CRYPTO-1 Stream Cipher



Principal Attack Technique

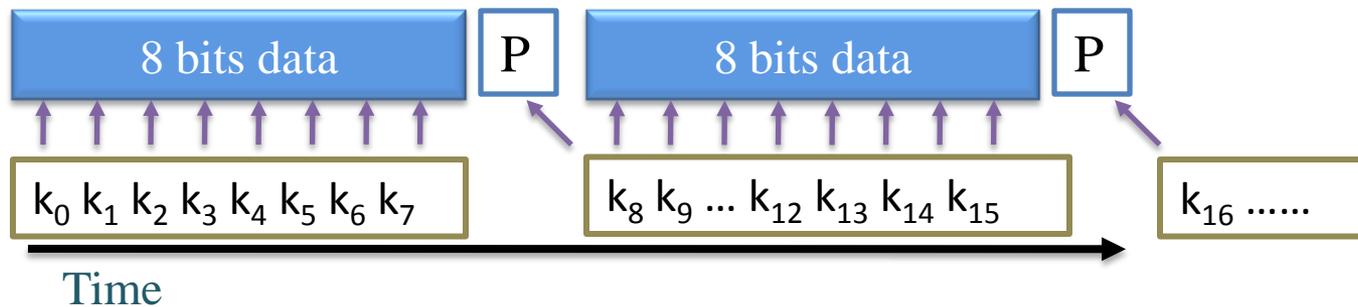
- Known-plaintext attack on stream cipher
 - ciphertext = plaintext XOR keystream
 - Ciphertext can be easily obtained via programmable reader or sniffer
 - If you know plaintext, then you know keystream
- Can recover internal state given enough keystream bits (plus enough computational power)

Main Vulnerabilities

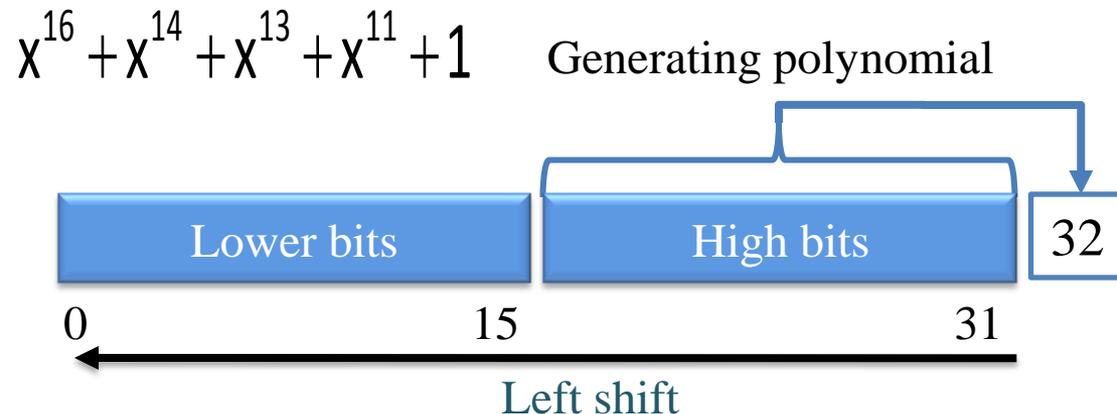
- CRYPTO-1's 48-bit key is way tooooooo short
 - Depending on which bits you have, the time to break can range from a few seconds to a few days
- Source of information leakage
 - Vulnerability in parity computation
 - Not enough entropy in nonce
 - Vulnerability in nonlinear filter function
 - Vulnerabilities in authentication protocol
 - Allows extremely efficient sniffer-based attacks

Parity and Nonce

- Parity against plaintext: Buy eight get one free

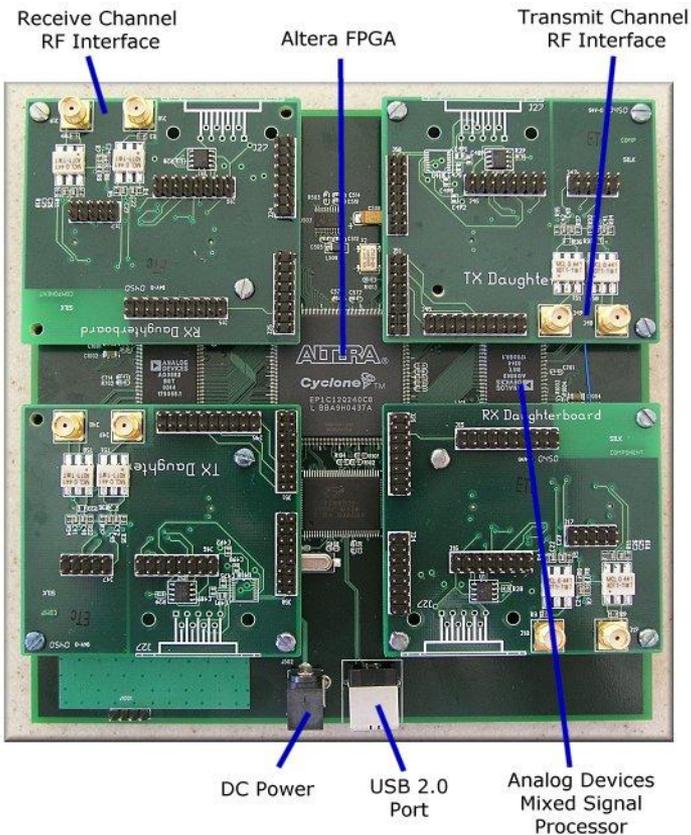


- 32-bit nonce function has only 16 bits of entropy



Equipment

Sniffer



PCD & PICC Emulator



Reader

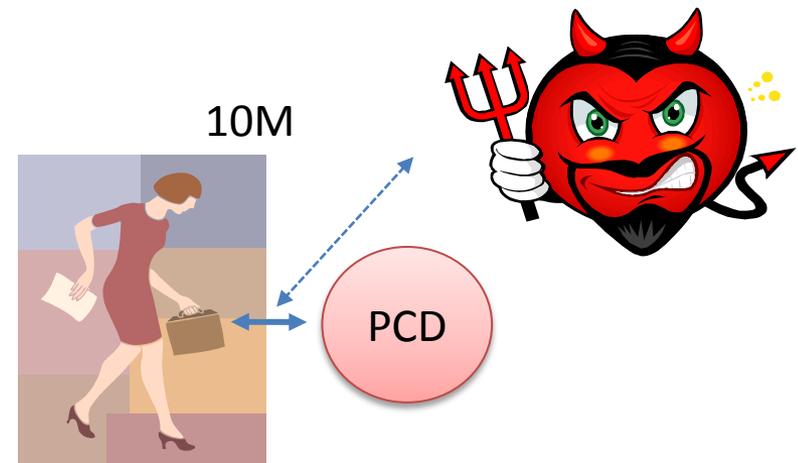
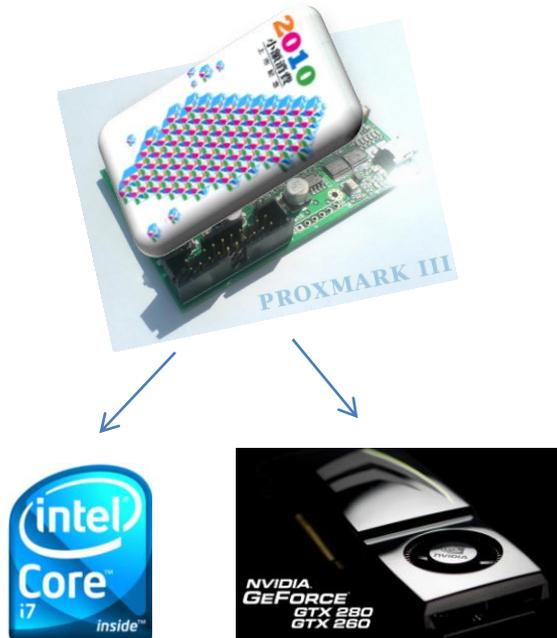


With MIFARE Classic chip

Attacks

PCD-based

Sniffer-based



Cost Comparison

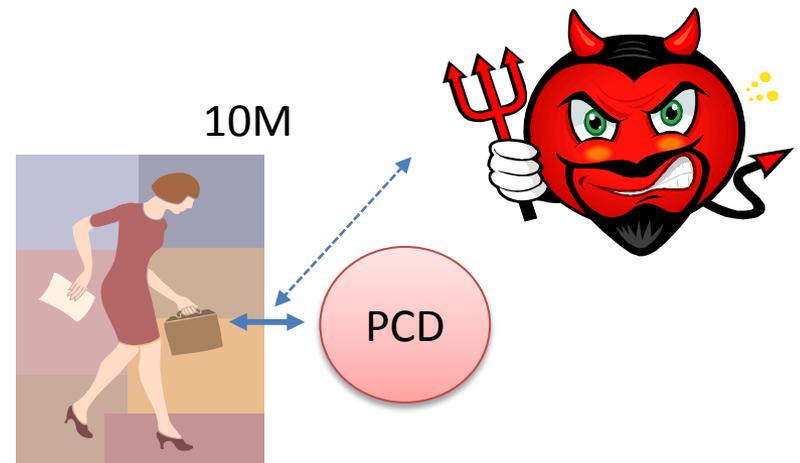
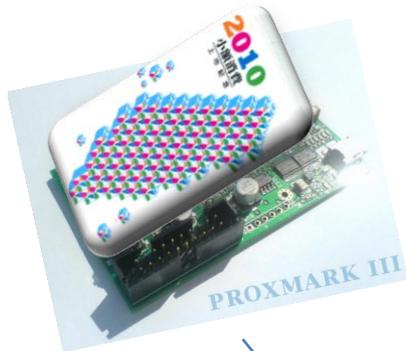


	PCD offline		Sniffer online
	First	Rest	Any
Platform	GPU	CPU	CPU
Devices	16	4	1
Time/per key	14 hour	1 hour	< 1 min

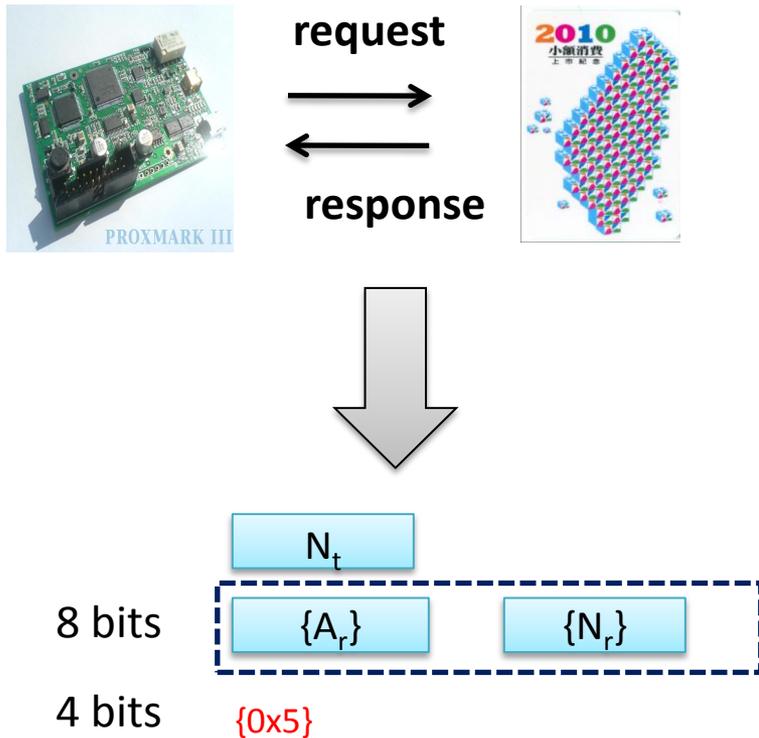
Attacks

PCD-based

Sniffer-based



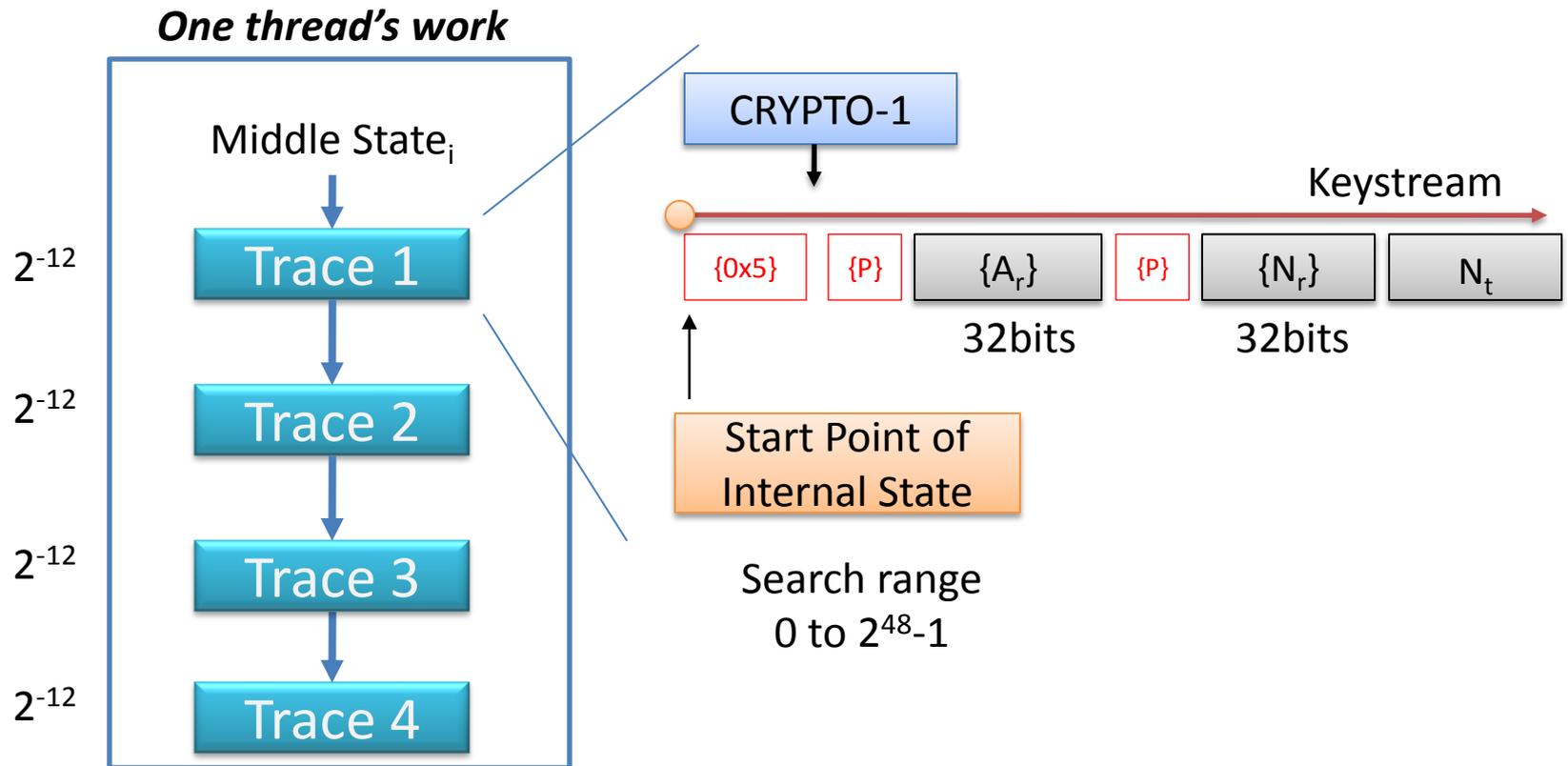
How to Obtain the First Key



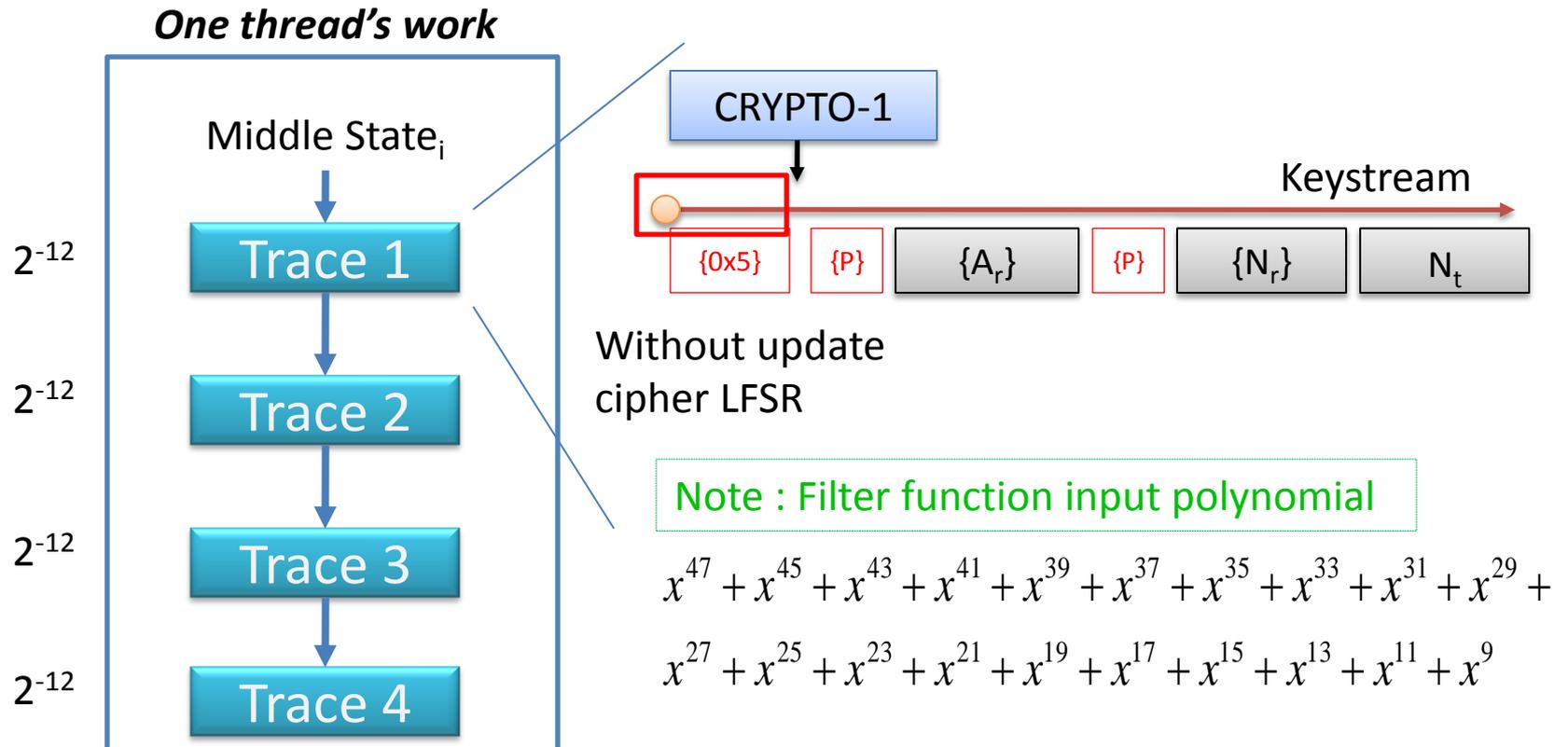
1. Keep requesting to authenticate
2. **4** to **6** traces
3. Brute-force search 2^{48} key space

PCD	PICC
6000f57b	
	f9105fce
{00000000} {00000000}	
{0} {0}	
	{5}
An error code trace	

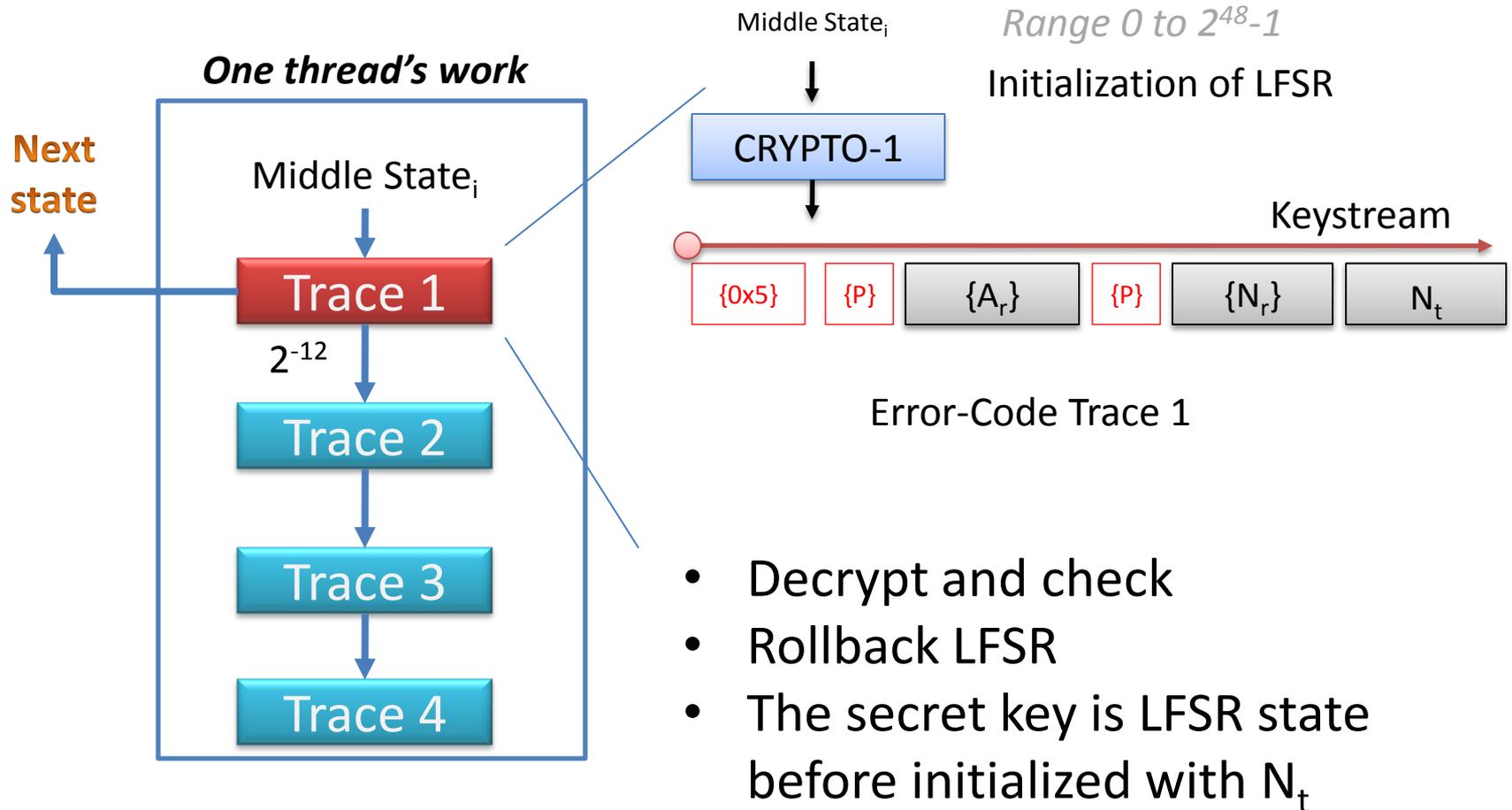
First Key by GPU Search



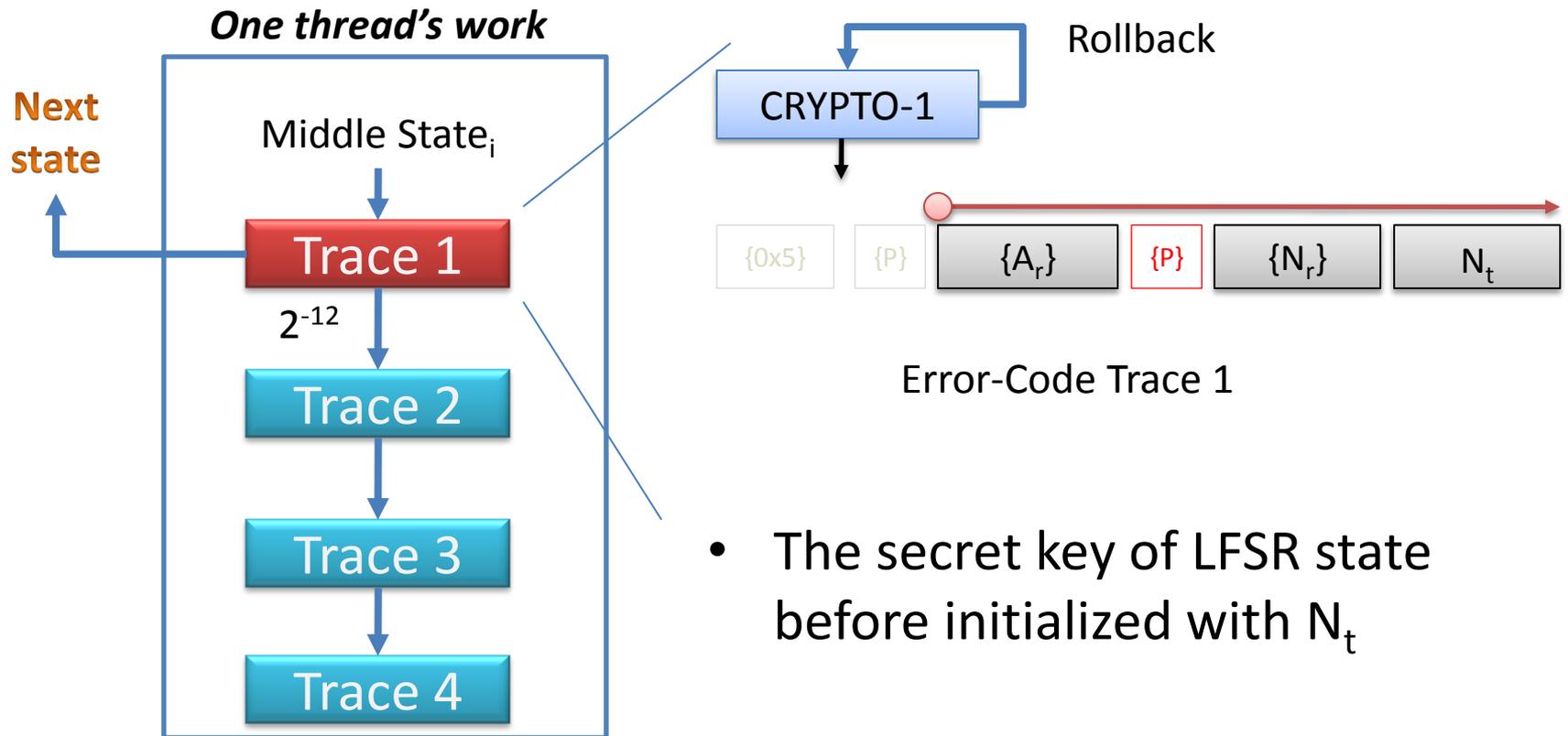
First Key by GPU Search



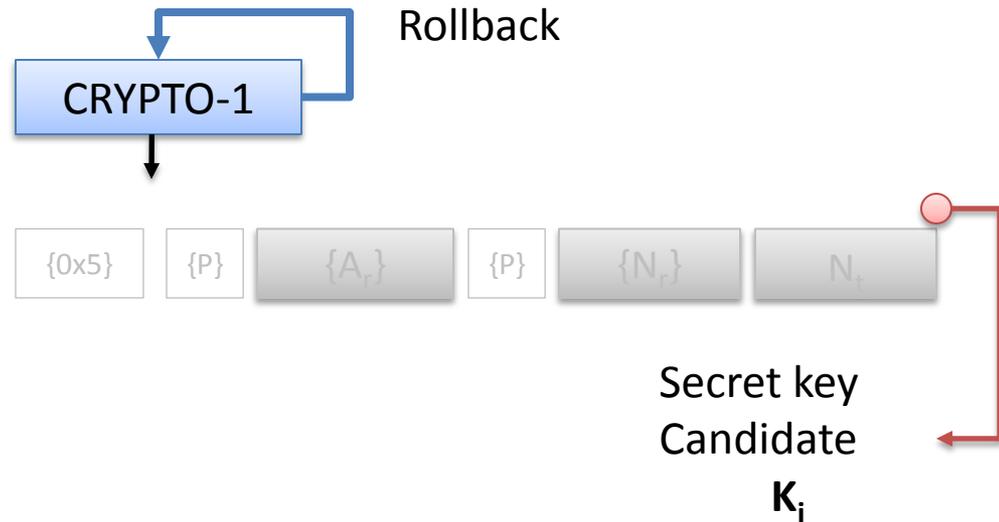
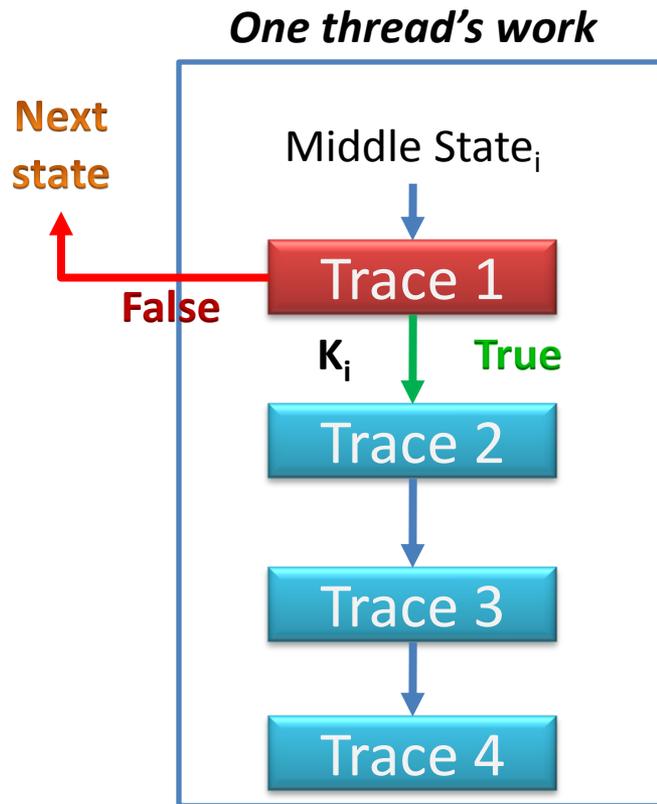
First Key by GPU Search



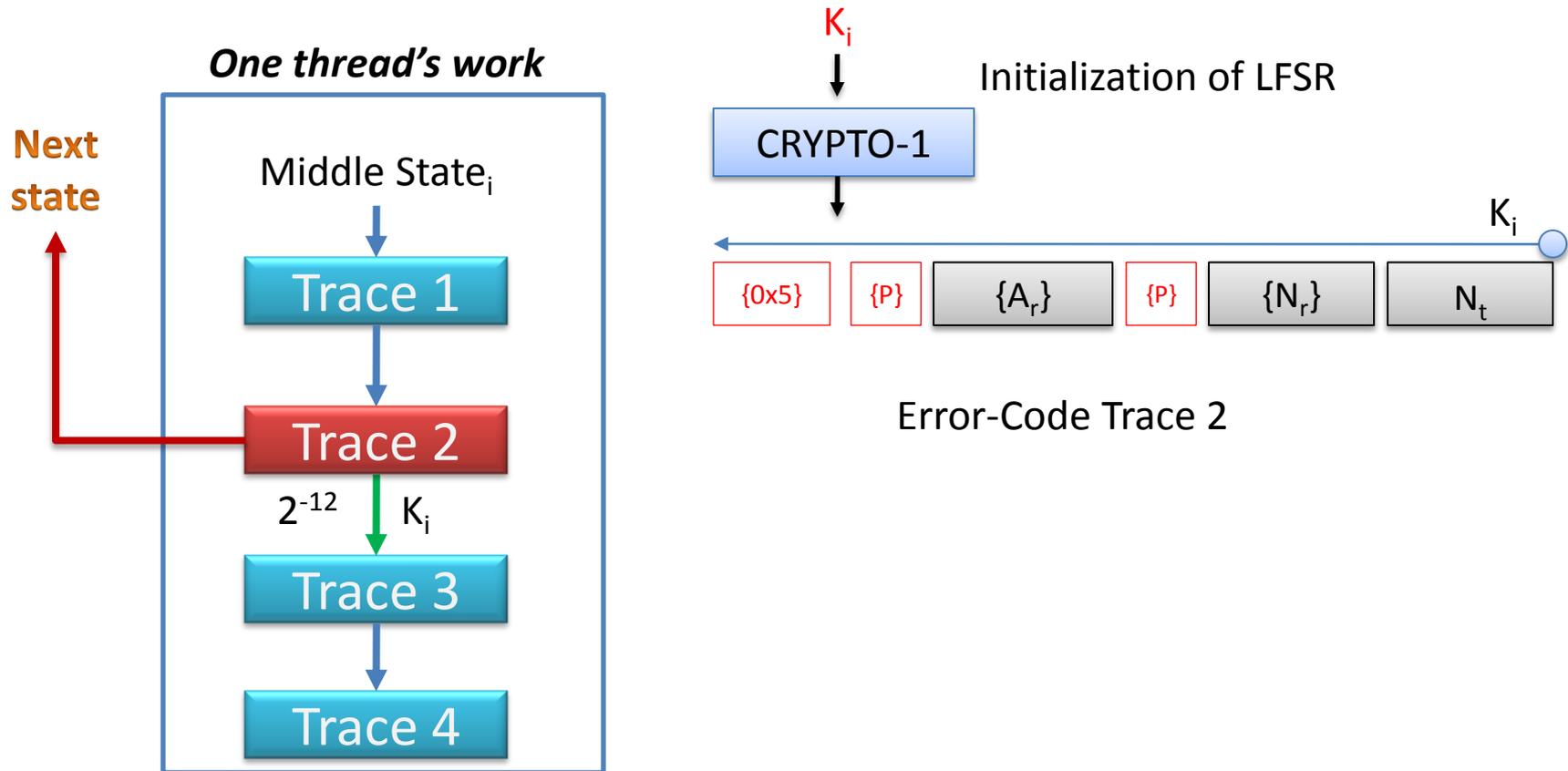
First Key by GPU Search



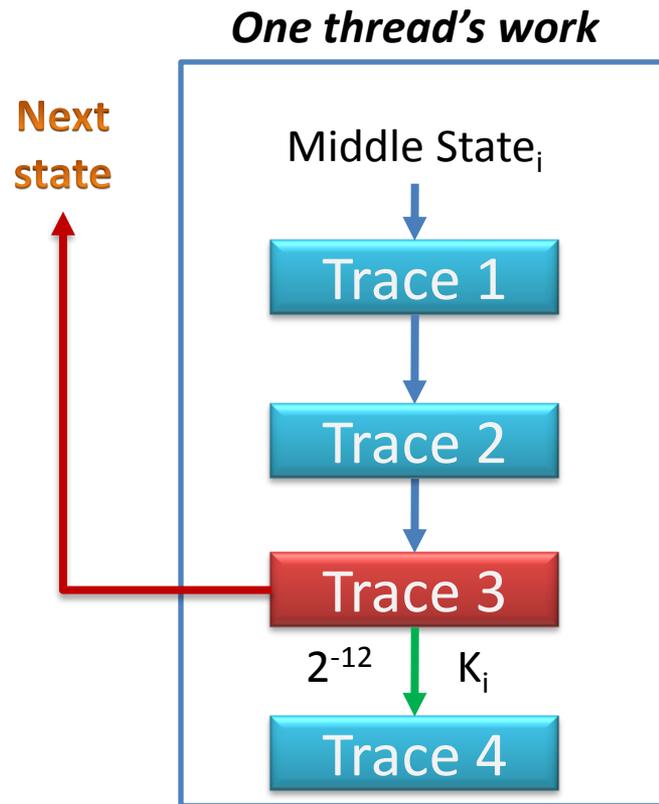
First Key by GPU Search



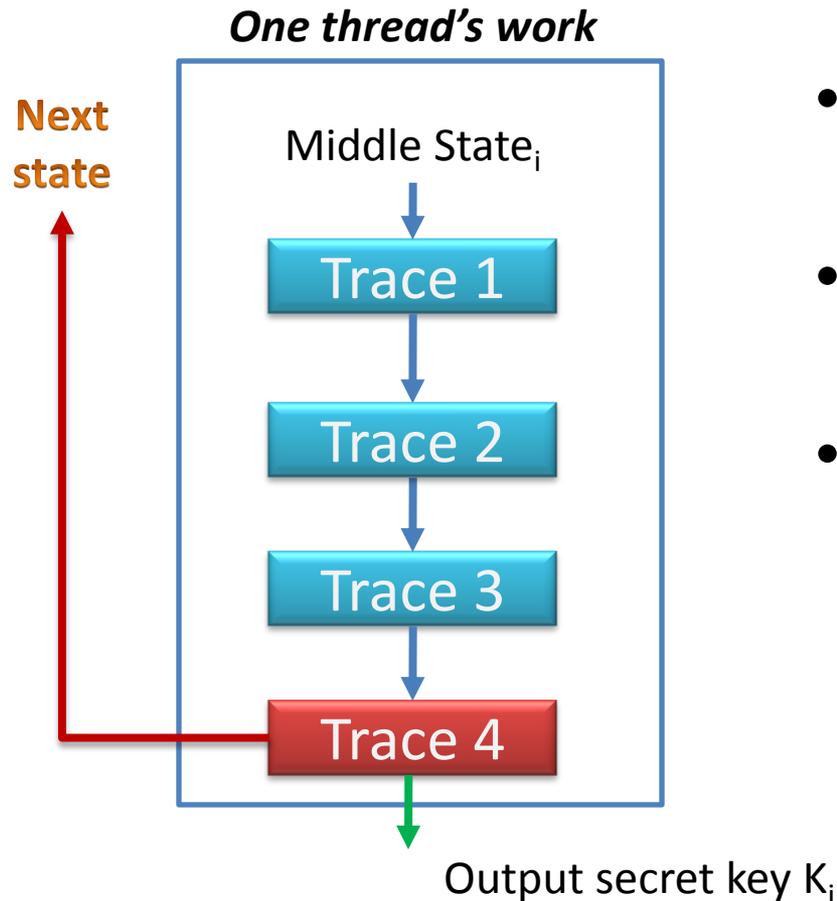
First Key by GPU Search



First Key by GPU Search



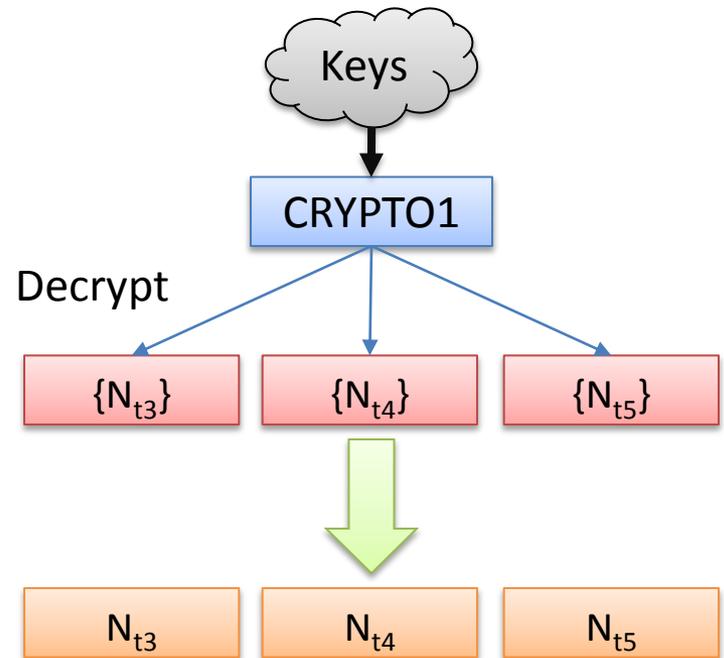
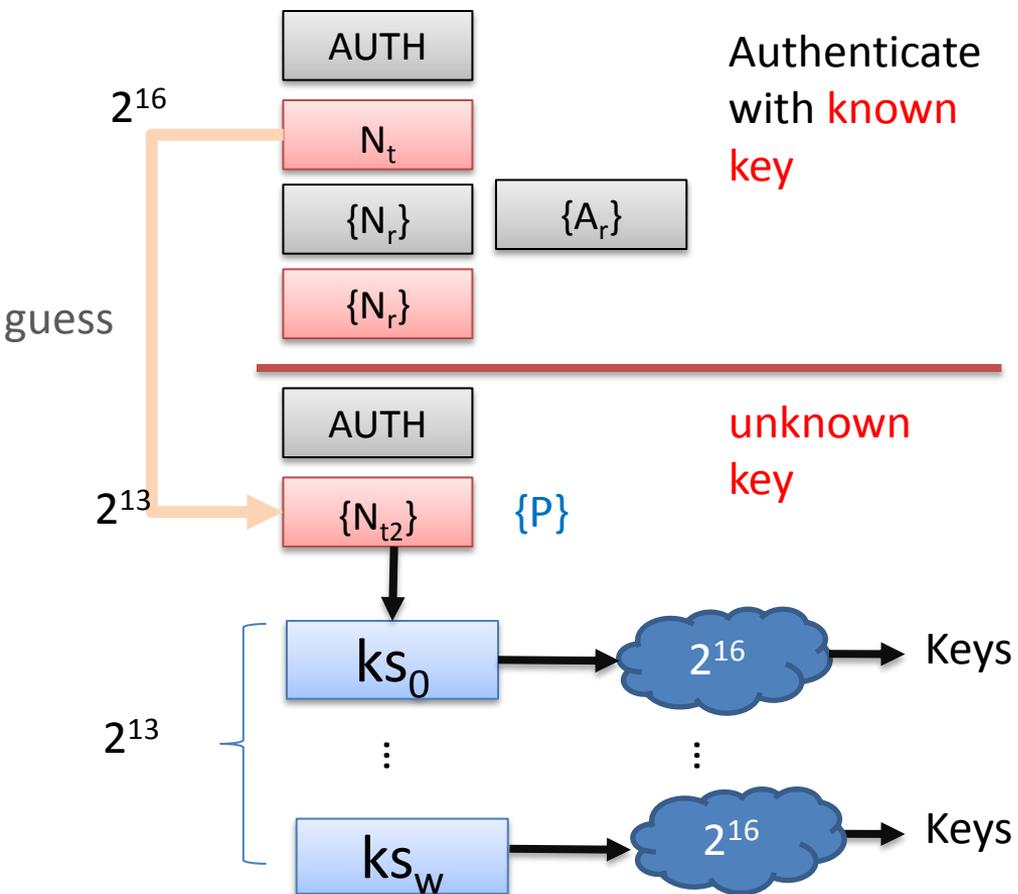
First Key by GPU Search



- Need at least four traces to decide unique secret key
- In practical, we run five or six traces
- **The speed of using four, five, and six traces is approximately same**

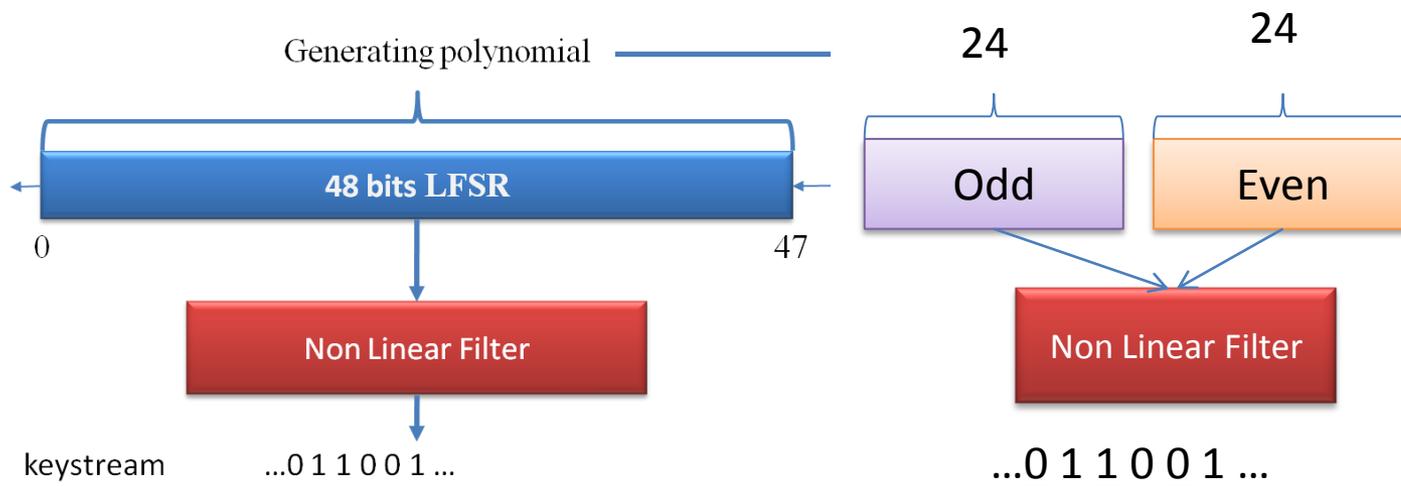
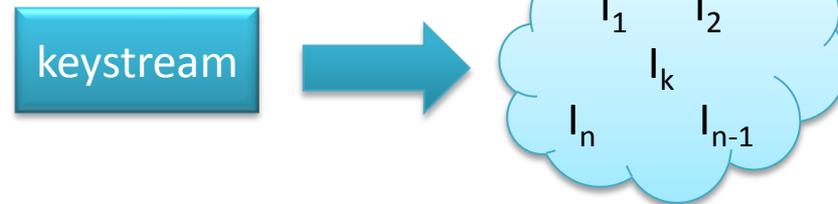
Getting Remaining Key

Nested authentication

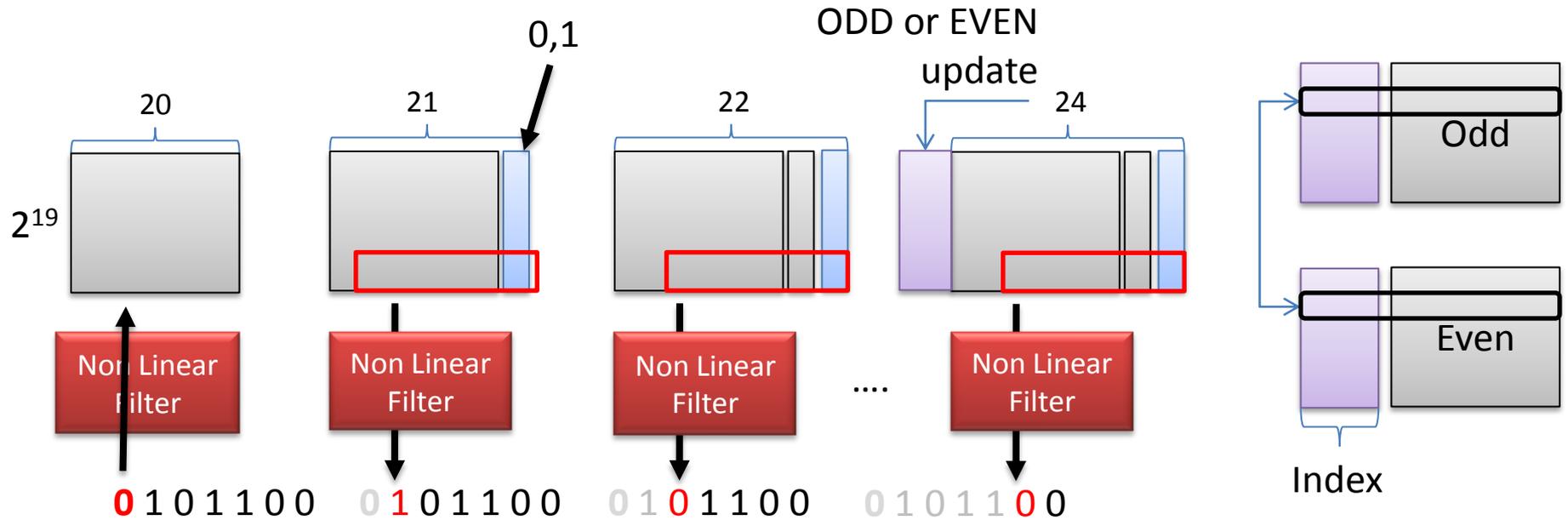


Inverting Filter Function

Garcia et al.
"Wirelessly pickpocketing a MIFARE Classic card."
In IEEE Symposium on Security and Privacy, 2009



A Time-memory Trade-off

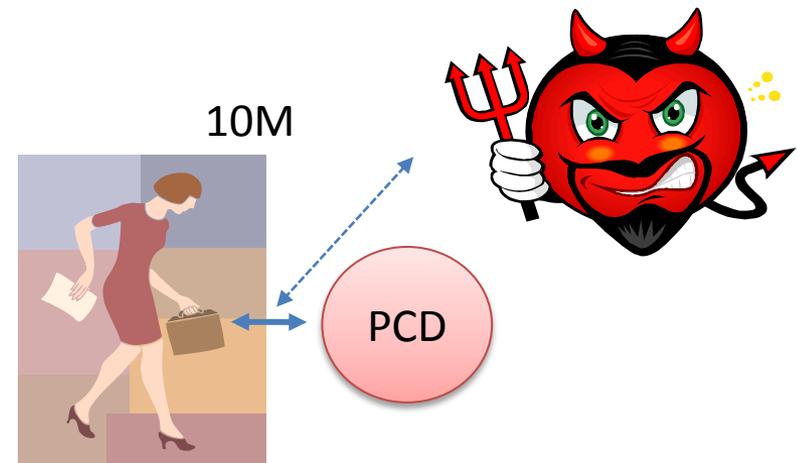
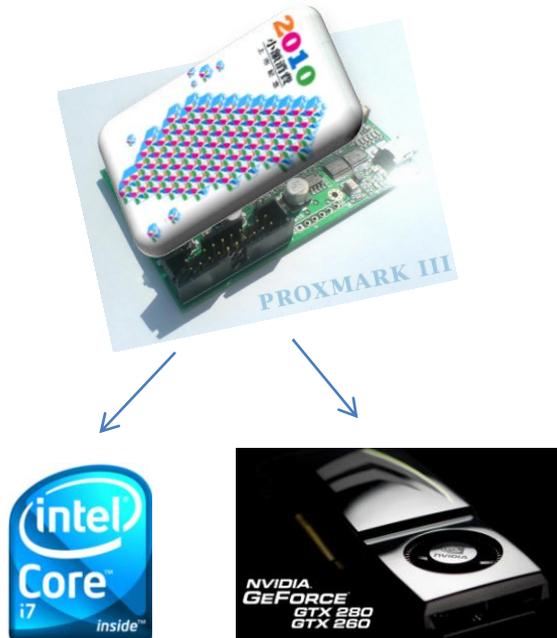


$$\begin{aligned}
 & X^{48} + X^{38} + X^{36} + X^{34} + X^{24} + X^6 + 1 \\
 +) & X^{43} + X^{39} + X^{33} + X^{31} + X^{29} + X^{23} + X^{21} + X^{19} + X^{13} + X^9 + X^7 + X^5
 \end{aligned}$$

Attacks

PCD-based

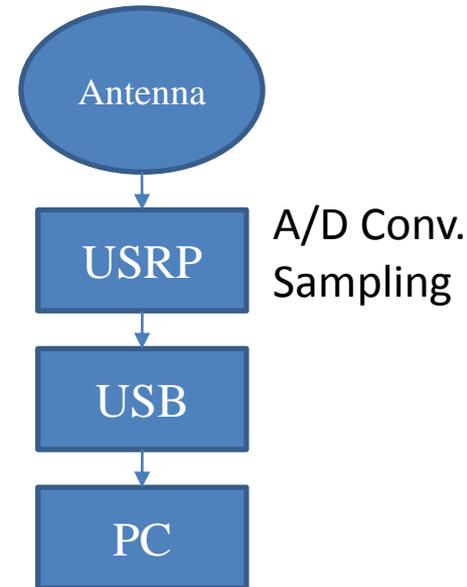
Sniffer-based



GNURadio-based Sniffer

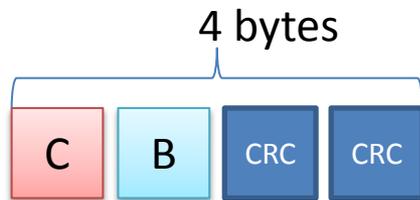
- Elements of the sniffer

1. A good antenna
2. USRP handles A/D and sampling
3. Transfer raw samples across USB
4. DSP on PC
 1. Demodulation
 2. Decoding
 3. Protocol analysis



Command Set

- Length of sequent transmission



Type	Bytes sequent	Function
V (INC, DEC, RES)	4-6-4	Change a value block
W (WRITE)	4-18	Write a block with 16 bytes data
A (AUTH)	4-8	Authenticate a sector by key A/B
R (READ)	4-next	Read a block

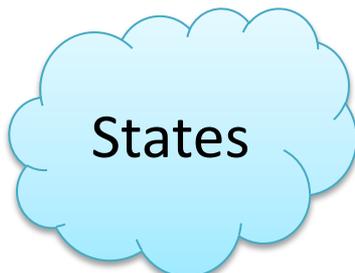
Inc/Res/Dec	Write	Authenticate	Read
{Inc/Dec/Res N} ₃₂	{Write N} ₃₂	Auth N ₃₂	{Read N} ₃₂
{ACK/NCK} ₄	{ACK/NCK} ₄	Nt ₃₂	{Data} ₁₄₄
{Value + CRC}₄₈	{Data CRC}₁₄₄	{Nr}₃₂ {Ar}₃₂	{Next Command}₃₂
{Transfer} ₃₂	{ACK/NCK} ₄	{At} ₃₂	
{ACK/NCK} ₄	{Next Command} ₃₂	{Next Command} ₃₂	
{Next Command} ₃₂			

Example One-way Trace

Anti-collision	
Auth 0x18 {NR} {AR} {Write 0x18} {write data}	6118e4fe 3edee7b0 3f307d3e 98c9b913 b1c903a22d1cc21b39d1502b894441473f00
{Auth 0x8} {NR} { AR } {DEC 0x8} {Value} {Transfer 0x8} {Read}	89be2cea 1433ad1452895e0c 8d02026d a2ef4ab078a9 84aaacec 5f815afa
{Auth 0x1a} {NR} { AR } {Write 0x1a} {Write Data}	fbf8c3d9 bcd863a91cf83b07 6fb38b89 72e4a262b284c235c7d054269d85e281d070
{Auth 0x10}	ff35fcc0

Example: WRITE Command

$$\begin{array}{|c|} \hline \text{a012cc82} \\ \hline \end{array} \oplus \begin{array}{|c|} \hline \text{98c9b913} \\ \hline \end{array} = \begin{array}{|c|} \hline \text{38db7591} \\ \hline \end{array}$$



CRYPTO-1

State_i



Anti-collision	
Auth 0x18	6118e4fe
N_t	
{N _r } ₃₂ {A _r } ₃₂	3edee7b0 3f307d3e
{Write 0x18} ₃₂	98c9b913
{ACK} ₄	
{write data} ₁₄₄	b1c903a2 2d1cc21b ...
{ACK} ₄	
{Auth 0x8} ₃₂	89be2cea

Decrypt trace to state_i

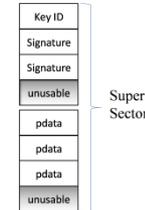
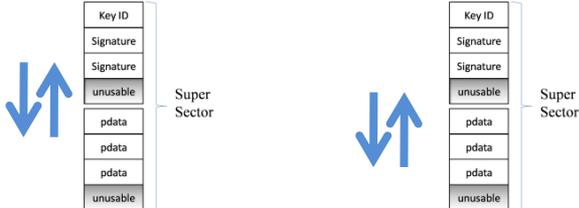
1. A_r is a MIFARE nonce

2. 0x89be2cea $\xrightarrow{?}$ 0x610865ee

Concluding Remarks: How to Fix MIFARE Classic?

- Under these attacks
MIFARE Classic is a **memory** card
- Need to defend against:
 1. Unauthorized content alteration
 2. Replay attack
 3. Clone attack
- Not unlike detecting counterfeit banknotes

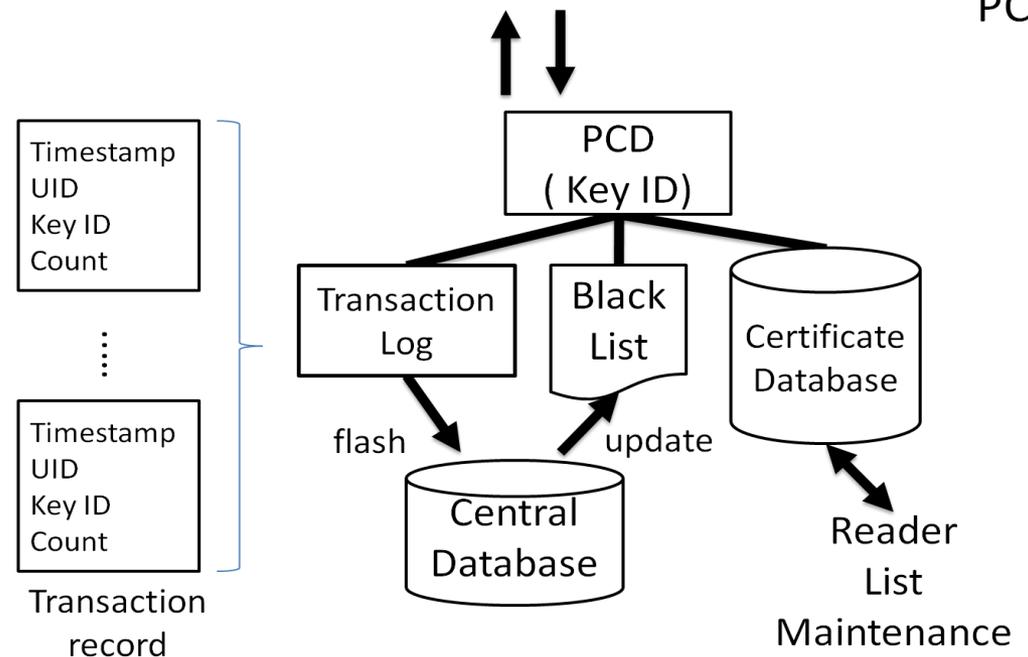
A Straightforward Defense Mechanism



Time synchronized



PCD



Timestamp
UID
Key ID
Count

⋮

Timestamp
UID
Key ID
Count

Transaction record

Value block

Key ID

Data block

Signature

Data block

Signature

unusable

Protecting data integrity using digital signature schemes

Example: TTS

Super Sector

Sector 0

UID

pdata

pdata

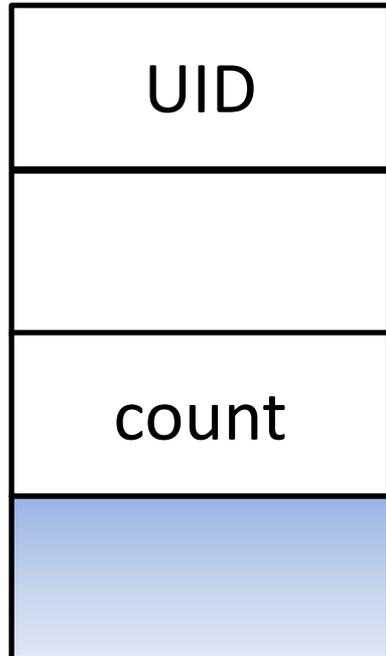
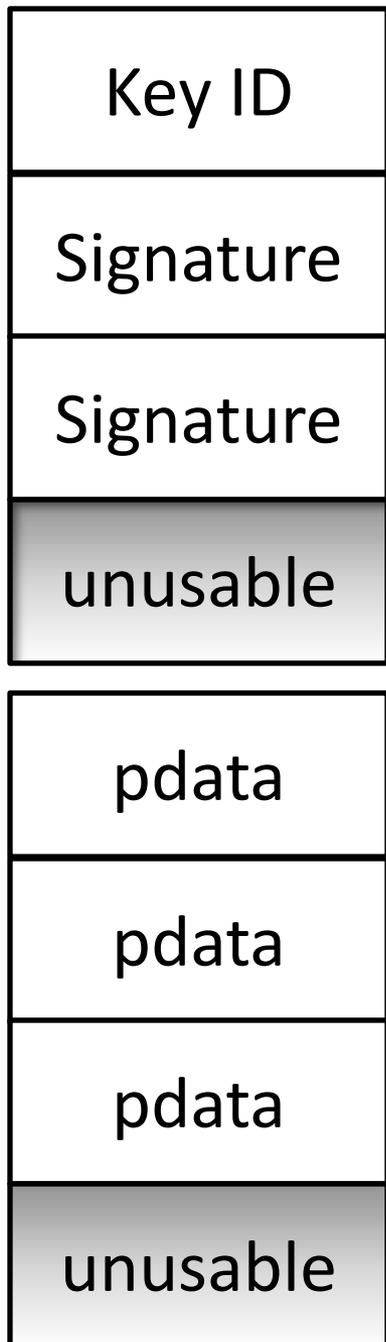
pdata

unusable

count

Data/Value block

PICC



If you are thinking to deploy MIFARE
Classic as a means of access control:
“Don’t.”

Thank you!

Questions or comments?